

Definability of linear equation systems over groups and rings

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A logic for polynomial time

Atserias, Bulatov, Dawar $\text{Slv}(\mathbf{G}) \notin \text{FP}+\text{C}$

Dawar, Grohe, Holm, Laubner $\text{FP}+\text{C} \not\leq \text{FP}+\text{rk} \leq \text{PTIME}$

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Matrix rank and linear equation systems

Fields $A \cdot x = b$ solvable iff $\text{rk}(A) = \text{rk}(A|b)$:

If $r = \text{rk}(A)$, then $a_1 \cdot c_1 + \dots + a_r \cdot c_r + a \cdot b = \mathbf{0}$

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Rings Many notions (linear dependence, McCoy, inner rank, ...),
unknown complexity, above characterisation fails

Groups Undefined

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Groups Undefined

Question: Is $\text{Slv}(\mathbf{G}) \in \text{FP}+\text{rk}$?

A systematic study of solvability

Inter-definability: \rightsquigarrow natural domain for Slv

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Theorem

k -ideal rings $\xrightarrow{\text{FP-red.}}$ cyclic groups of prime power order.

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Theorem

Normal form for $\text{FO} + \text{slv}_F$.

Inter-definability: a natural class for solvability

$\text{Slv}(\mathbf{CG})$: Cyclic groups (\mathbb{Z}_p^e)

$\text{Slv}(\mathbf{I}_k\mathbf{R})$: k -gen. ideal rings ($I \trianglelefteq R \Rightarrow I = \pi_1 R + \cdots + \pi_k R$)

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R **local** iff $R \setminus R^* \trianglelefteq R$

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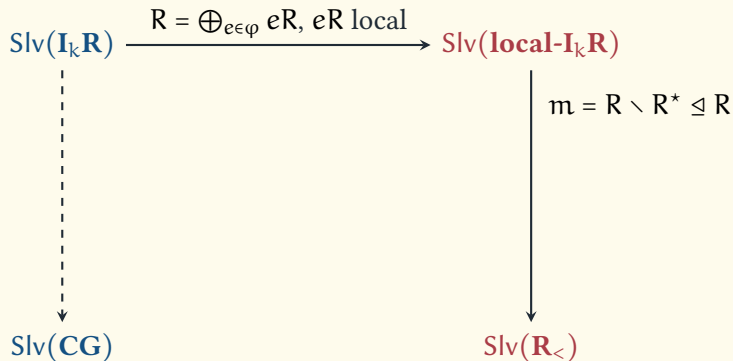
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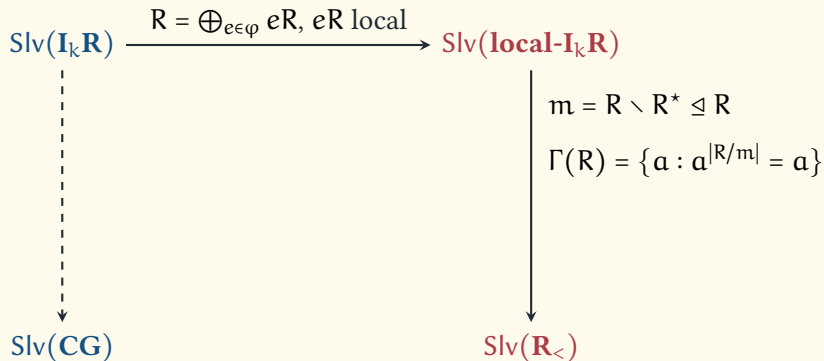
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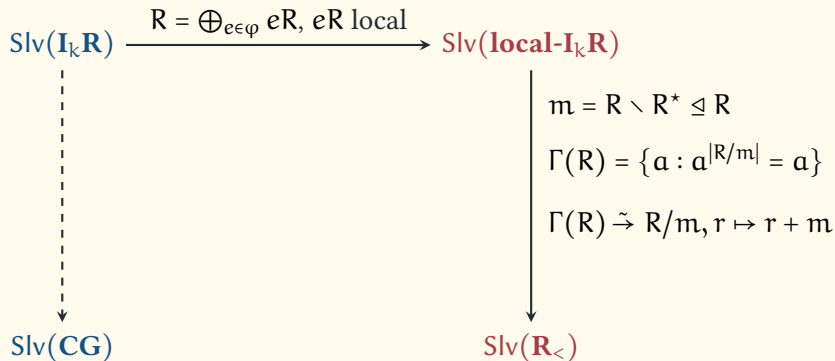
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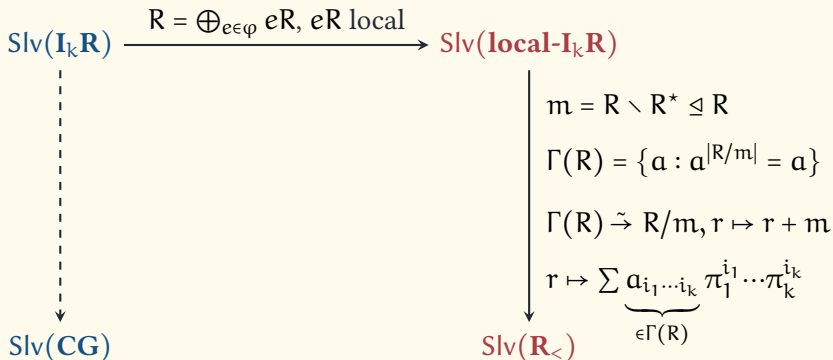
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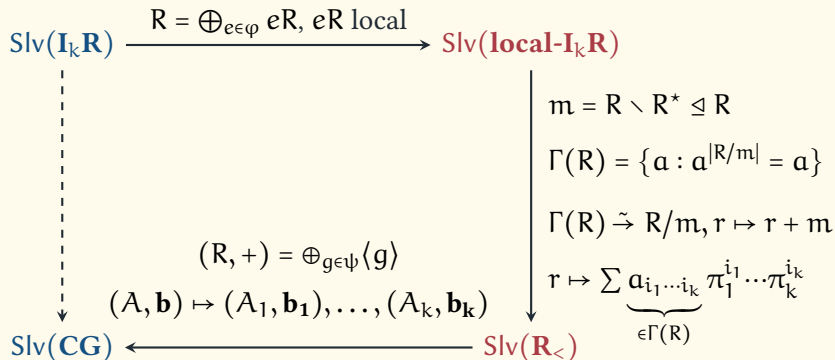
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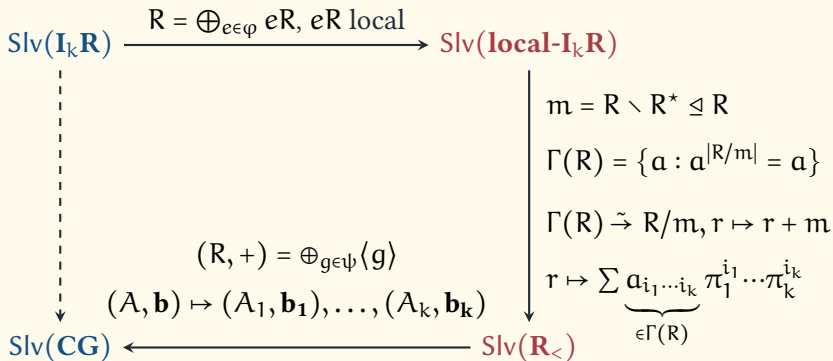
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Theorem $\text{Slv}(\mathbf{I}_k\mathbf{R}) \leq_{\text{FP}}^{\text{T}} \text{Slv}(\mathbf{CG})$

Intra-definability: solvability as a logical operator

$$\text{slv}(\bar{x}, \bar{y}, \bar{r}_i). \left[\varphi_M(\bar{x}, \bar{y}, \bar{r}), \varphi_b(\bar{x}, \bar{r}), \underbrace{(\varphi_R, \varphi_+, \varphi_.)}_{\text{finite ring}}(\bar{r}_1, \bar{r}_2, \bar{r}_3) \right]$$

coefficient matrix solution vector finite ring

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FO+slv : First-order logic closed under solvability quantifier

FO+slv_F : Solvability quantifier over a fixed finite field F

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Theorem

Every $\text{FO} + \text{slv}_F$ -formula equivalent to an $\text{FO} + \text{slv}_F$ -formula

$$\text{slv}(\bar{x}, \bar{y}).[\varphi_M(\bar{x}, \bar{y}), \mathbf{1}], \text{ with } \varphi_M \text{ quantifier-free.}$$

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Proof illustration: (negation)

$$\neg \text{slv}(\bar{x}, \bar{y}).[\varphi, \mathbf{1}]$$

$$\text{Non-solvability} \equiv \neg \exists \mathbf{x} : \mathbf{M}\mathbf{x} = \mathbf{b} \stackrel{?}{\equiv} \exists \mathbf{y} : \mathbf{M}'\mathbf{y} = \mathbf{b}' \equiv \text{Solvability}$$

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Gaussian elimination implies:

$$\neg \exists \mathbf{x} : \mathbf{M}\mathbf{x} = \mathbf{b} \equiv \exists \mathbf{y} : \mathbf{y}(\mathbf{M}|\mathbf{b}) = (0, \dots, 0|1).$$

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Proof illustration: (conjunction)

$$\text{slv}(\bar{x}, \bar{y}).[\varphi, \mathbf{1}] \wedge \text{slv}(\bar{x}, \bar{y}).[\psi, \mathbf{1}]$$

$$\boxed{\varphi} \cdot \mathbf{v}_y = \boxed{\begin{matrix} 1 \\ \vdots \\ 1 \end{matrix}}$$

$$\boxed{\psi} \cdot \mathbf{v}_y = \boxed{\begin{matrix} 1 \\ \vdots \\ 1 \end{matrix}}$$

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φ	$\cdot \mathbf{v}_y =$	$\begin{matrix} 1 \\ \vdots \\ 1 \end{matrix}$	\rightsquigarrow	<table style="border-collapse: collapse; margin: auto;"><tr><td style="padding: 5px;">φ</td><td style="border-left: 1px solid black; padding: 5px;">$\mathbf{0}$</td></tr><tr><td colspan="2" style="border-top: 1px solid black;"></td></tr><tr><td style="padding: 5px;">$\mathbf{0}$</td><td style="border-left: 1px solid black; padding: 5px;">ψ</td></tr></table>	φ	$\mathbf{0}$			$\mathbf{0}$	ψ	$\cdot \mathbf{v}_{yy} =$	$\begin{matrix} 1 \\ \vdots \\ 1 \end{matrix}$
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Proof illustration: (universal quantification)

$$\forall z (\text{slv}(\bar{x}, \bar{y}) \cdot [\varphi(\bar{x}, \bar{y}, z), \mathbf{1}])$$

$$\boxed{\varphi(z_1)} \cdot \mathbf{v}_y = \boxed{\begin{matrix} 1 \\ \vdots \\ 1 \end{matrix}}$$

\vdots

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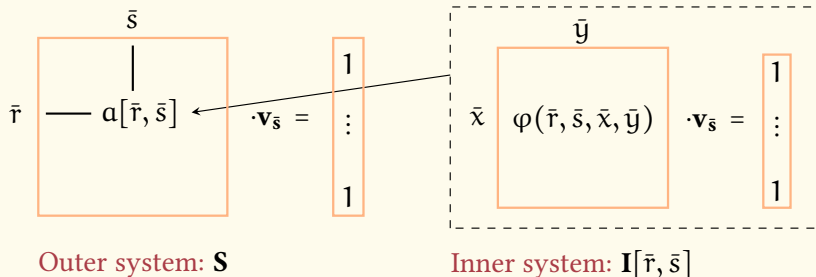
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Proof illustration: (nesting of solvability)

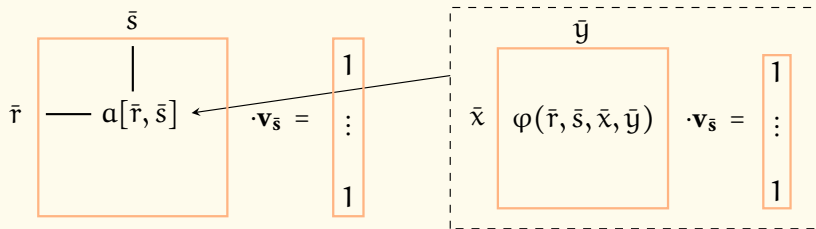
$$\text{slv}(\bar{r}, \bar{s}).[\text{slv}(\bar{x}, \bar{y}).[\varphi(\bar{r}, \bar{s}, \bar{x}, \bar{y}), \mathbf{1}], \mathbf{1}]$$



Intra-definability: solvability as a logical operator

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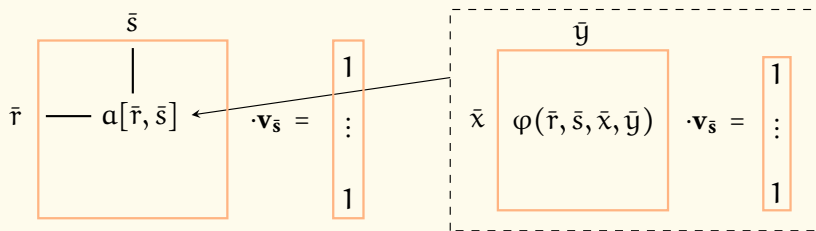
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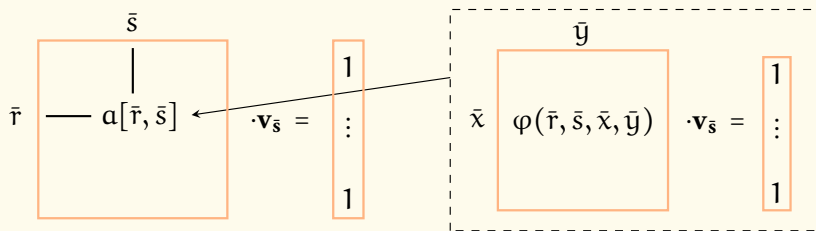


$$\text{For } \bar{r}: \underbrace{\sum_{\bar{s}} a[\bar{r}, \bar{s}] \cdot v_{\bar{s}}}_{= 1} = 1$$

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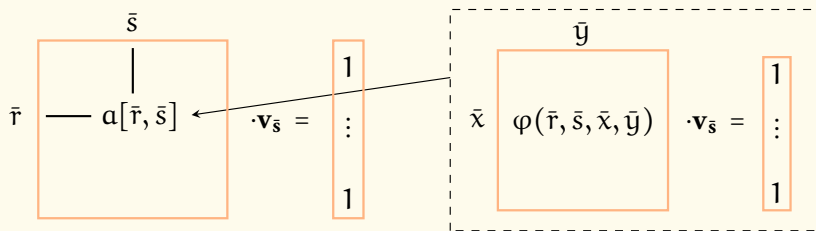
$$\text{For } \bar{r}: \underbrace{\sum_{\bar{s}} a[\bar{r}, \bar{s}] \cdot v_{\bar{s}}}_{\Downarrow} = 1$$

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Consistency conditions:

$$v[\bar{r}, \bar{s}] = 1 \Rightarrow a[\bar{r}, \bar{s}] = 1$$

$$v[\bar{r}, \bar{s}] \neq v[\bar{r}', \bar{s}] \Rightarrow a[\bar{r}, \bar{s}] \neq a[\bar{r}', \bar{s}]$$

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How to formalise: “If $v = 1$ then $A \cdot x = 1$ solvable”

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How to formalise: “If $v = 1$ then $A \cdot x = 1$ solvable”

$$\boxed{A} \begin{array}{c} -v + 1 \\ \vdots \\ -v + 1 \end{array} \cdot \mathbf{x} = \boxed{\begin{array}{c} 1 \\ \vdots \\ 1 \end{array}}$$

Conclusion and outlook

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Outlook: Permutation group membership (GM)

Given: Permutations π_1, \dots, π_k and π on a set A

Question: Is $\pi \in \langle \pi_1, \dots, \pi_k \rangle \leq S_A$?

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$$\text{Slv}(\mathbf{D}) \xrightarrow{\text{FO-reduction}} \text{GM} \quad (\text{Cayley's theorem})$$